Optimal Cable Laying Across an Earthquake Fault Line Considering Elliptical Failures

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Abstract—Whether it be a road, railway track, pipeline or fibre-optic cable, achieving a low probability of “cable” damage is important for the proper operation of modern infrastructure. This paper considers the fundamental problem of how to connect two points with a cable that crosses an earthquake fault line. We first develop a model, under certain general assumptions, for the cable break probability. Then we formulate a multi-objective optimization problem, with cable cost and probability of cable break being the two objectives. For two important and meaningful sets of cable shape alternatives, the Pareto front is determined for these two objectives. All the analytical results are verified by simulations.

Index Terms—Cabling, cost minimization, earthquake, disaster, survivability, probability, pipeline, infrastructure.

ACRONYMS AND ABBREVIATIONS:

<table>
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<tr>
<th>Abbreviation</th>
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<tr>
<td>CBP</td>
<td>Cable Break Probability</td>
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<tr>
<td>CCB</td>
<td>Conditional Cable Break</td>
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<tr>
<td>CSP</td>
<td>Cable Survival Probability</td>
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<td>EFL</td>
<td>Earthquake Fault Line</td>
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<td>HRA</td>
<td>Hook with Right Angles</td>
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<td>TCS</td>
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NOTATIONS:

\[ S = (x_S, y_S) \]

A point \( S \) in \( \mathbb{R}^2 \) with coordinates \((x_S, y_S)\). Points in \( \mathbb{R}^2 \) are denoted by uppercase letters, while the coordinates of such points are denoted by lowercase letters.

\( A, B \) two endpoints of the EFL

\( D \) set of points on the EFL

\( \hat{D} \) length of the EFL

\( N_1, N_2 \) two points to be connected

\( e_1, e_2 \) coefficients of the linear function that describes the straight line \( \overrightarrow{N_1N_2} \)

\( \hat{d} \) distance between \( N_1 \) and \( N_2 \)

\( L \) set of points on the cable

\( c \) cable cost per unit length

\( \varepsilon \) cost of the cable

\( \varepsilon \) eccentricity of the elliptical isoseismals

\( \sqrt{1 - e^2} \) equivalent distance from \( S \) to \( M \)

\( d(S) \) equivalent distance from \( S \) to the cable

\( P(\text{break}) \) CBP function

\( \lambda \) a parameter used in the numerical examples representing the exponential decay rate of the conditional cable break probability

\( \overrightarrow{S_1S_2} \) line that contains \( S_1 \) and \( S_2 \)

\( \overrightarrow{S_1S_2} \) line segment between \( S_1 \) and \( S_2 \)

\( |S_1S_2| \) length of \( \overrightarrow{S_1S_2} \)

\( \Delta_T \) length of middle segment of a TCS cable

\( \Delta_H \) length of last segment of a HRA cable

\( d_1(S) - d_3(S) \) three candidates for \( d(S) \) under the straight line cable shape

\( d_{T1}(S) - d_{T5}(S) \) five candidates for \( d(S) \) under TCS

\( d_{H1}(S) - d_{H3}(S) \) three candidates for \( d(S) \) under HRA

I. INTRODUCTION

In many applications, engineers are faced with a challenge of how to lay an optical fibre cable, a fuel or gas pipeline, electric power line, road or railway track between two points across an earthquake fault zone. Connecting the two points by a straight line will minimize cost but may increase the risk of damage or break in a case of an earthquake. In practice, engineers lay cables using a software package e.g.
MakaiPlan [1] and may consider potential earthquake hazards by introducing ad hoc adjustments. In this paper, we develop mathematical expressions and related numerical results for cable cost and risk of break that will provide new insights into how to lay a single cable between two points across an Earthquake Fault Line (EFL) assuming that earthquake epicenters only occur on the EFL.

Modern infrastructures have been adversely affected by a series of earthquakes with grave and costly consequences over the last century. Several examples of such earthquakes follow. In 1971, an earthquake in San Fernando caused about 1,400 breaks in various piping systems, which resulted in temporal disruption of water, gas and sewage services of the city [2]. In 1987, an earthquake in Colombia-Ecuador border region destroyed or badly damaged about 27 km of the oil pipeline in Ecuador, which may represented the largest single pipeline loss in history [3]. An earthquake of 7.1-magnitude on 26th December 2006 just south of Taiwan caused unprecedented damage to seven submarine communication cables. This incident caused severe disruption of Internet and phone services in south east Asia for several weeks [4]. Similar multiple-cable breaks happened twice in the Mediterranean sea in 2008, where the undersea cables connecting Europe, North Africa and Asia are concentrated in a small area [5]. A massive earthquake of 9.0-magnitude followed by a tsunami occurred near the east coast of Honshu, Japan on 11th March 2011, which destroyed or damaged at least 2,126 roads and 26 railways [6]. Many other strong and destructive earthquakes have occurred in the past and are likely to occur in the future.

There have been many research publications aiming to understand the effect of earthquakes and improve the survivability of modern infrastructures [7], [8]. Several researchers have studied certain past earthquakes to evaluate their potential damage on current infrastructures. For example, [9] evaluated the seismic risk of the water supply pipeline system in Los Angeles. [10], [11] investigated the potential damage to current natural gas or water pipelines in Taichung City, Taiwan from the 1999 Chi-Chi earthquake (also called Ji-Ji earthquake). Other researchers have focused on modeling and analyses of the vulnerability of infrastructure systems. The work presented in [12] considered power and gas pipeline systems, and proposed a methodological framework for analyzing the vulnerability of interdependent infrastructure systems. There is a method based on the identification of the parameters developed in [13] to assess the seismic vulnerability of water pipeline network. The study presented in [14] developed a Geographic Information Systems based probabilistic reliability model to assess seismic vulnerability of European gas and electricity networks in terms of various performance measures.

It is hard to develop a realistic disaster model, since disaster can happen anytime anywhere and the level of impact can vary as well [15]. Therefore, many researchers use computational geometric methods to model disasters and evaluate network resilience [16], [17]. For example, [18], [19] proposed a grid partition-based scheme to assess the network vulnerability under random region failures. In [20], the effects of disasters were considered as random line-cuts. The works presented in [21]–[24] attempted to identify the vulnerable points in geographic networks by finding the location of the worst-case failure (a line or disk cut) using certain network connectivity measures. In the research presented in [25]–[28], it was assumed that a disaster occurs as a randomly placed circular disk of a particular radius. This is known as the circular disk failure model. In [29], [30], circular disk failures were also considered in small networks to solve a problem of cost minimization subject to meeting a survival probability constraint.

Although the circular disk failure model [25]–[30] is popular, its generalization, the elliptical failure model is more suitable for many especially strong earthquakes. The elliptical failure model properly incorporates the direction dependent effects of earthquakes. Observations from many past earthquakes, such as the Charlevoix-Kamouraska earthquake (1925) [31], New Zealand earthquake (1929) [32], the Koyna earthquake (1967) [33], the Tonghai earthquake (1970) [34], the Tangshan earthquake (1976) [35], and the Kocaeli earthquake (1999) [36], have shown that the earthquake effects are indeed direction dependent and the shape of the isoseismals is elliptical. Note that, although an earthquake epicenter is frequently visualized as a point, the fault rupture inside the Earth’s crust that causes the earthquake, in fact, occurs along a 3D plane and often propagates as an intermittent series of fault ruptures or slips within the ruptured zone. Along the parallel direction of fault ruptures, the earthquake effect tends to accumulate and becomes growingly significant as the intermittent series of fault ruptures occurs and propagates. In contrast, along the normal direction of fault ruptures, the accumulation of earthquake effect is less significant. This fault rupture mechanism leads to the common observation of elliptical isoseismals in many past earthquakes [37]. When the earthquake magnitude is small, the fault rupture zone is small and the earthquake epicenter may be well approximated as a point (i.e., no intermittent series of fault ruptures, but just a single event of fault rupture at a point). Therefore, the direction dependent effect is not significant and the elliptical isoseismals gradually converge to circular isoseismals [37]. When developing a general earthquake effect model without reference to specific fault systems or specific sites, an elliptical failure model is able to properly capture the direction-dependent effects of earthquakes. Clearly, the circular disk failure is a special case of the elliptical failure when the eccentricity of the ellipse is equal to 0, and therefore, all the results of this paper are also applicable to the circular disk failure model used in [25]–[30].

In addition to the elliptical failure assumption, we also assume that the conditional probability of a cable break (for a given epicenter) is a generally decaying function of a certain measure of the cable position and the location of the epicenter.

In addition to the above mentioned publications, there were many others (e.g. [38]–[47] and references therein) that have studied network survivability subject to one or more link or node failures. However, not many publications have focused on minimizing the risk of an individual link or cable failure.
This is the topic of the present paper.

There are many pairs of cities in the world that are directly connected by at least one cable that crosses at least one earthquake fault. Examples include: Lanzhou and Chengdu, San Jose and Los Angeles [48]–[50], etc.

The unique feature of this work is that we aim to optimize the shape of the single cable and to derive convenient expressions obtained by probabilistic analysis. This has not been done by any of the previous publications. We note that in [29], [30] certain cable shape scenarios have been considered for multiple cables between two nodes, but the results of [29], [30] are not applicable to the case we consider here of laying a single cable. In many cases, infrastructure vendors focus on the problem of laying a single cable and it is therefore an important problem on its own. A key performance measure in this paper is the cable break probability (CBP). It represents the probability that a cable breaks following an earthquake.

A related concept is the cable survival probability (CSP), which is one minus the CBP. Meeting cable survivability requirement is an important design objective in addition to cost. We consider the problem of how to minimize the cost of laying a cable across an EFL, to maintain a required cable survivability probability in an earthquake event, or to minimize the CBP for a given budget. To this end, we formulate it as a multi-objective optimization problem that considers cable cost and CBP as the two objectives. We consider two important and meaningful sets of cable shape alternatives, and provide Pareto fronts for the two objectives. For tractability, to be able to obtain convenient expressions for determining the Pareto fronts, we rely on simplifications of a far more complex real world problem. Nevertheless, our solutions provide certain insight and useful guidance for cable laying. Furthermore, although the work here is written in the context of cabling, it has many other applications, including designs of fuel or gas pipelines, roads, railway tracks and power lines.

The remainder of the paper is organized as follows. In Section II, we describe the problem. In Section III, we provide a brief description and intuitive justification of the cable shape alternatives. In Section IV, we consider the cases of straight-line cables which provide insight into how to lay cables with low CBP. In Sections V and VI, we study two important alternatives of cable shapes and obtain expressions for the cost and CBP. After establishing the relevant monotonicity properties of the convenient probabilistic expressions, we provide Pareto fronts for cost and survival probabilities of the two alternatives. The paper is concluded in Section VII.

II. PRELIMINARIES AND PROBLEM DESCRIPTION

As discussed, we consider tradeoff between the two objectives: cable cost and CBP. To obtain the Pareto front for these two objectives, we will provide proofs for their monotonicity. Having the Pareto front enables us to solve two single-objective optimization problems of minimizing cost subject to a constraint on the CBP, and minimizing CBP subject to a constraint on the cable cost, in a case of an earthquake.

The problem is defined in a 2-dimensional Euclidean space \( \mathbb{R}^2 \). We assume that only one destructive earthquake occurs at one time. This is justified by the low frequency of earthquakes and by the relatively short time cables are being repaired. Let \( \mathcal{D} \subset \mathbb{R}^2 \) be a set of points on the EFL where an earthquake epicenter may occur. A coordinate system is constructed so that the origin is at the midpoint between the two endpoints of the EFL: \( A \) and \( B \). We assume for simplicity that the set \( \mathcal{D} \) is a segment of a straight line \( y = 0 \). Let \( \hat{D} \) be the length of the EFL. Thus \( A \) and \( B \) are positioned at the \( x \)-axis points \((-\hat{D}/2, 0)\) and \((\hat{D}/2, 0)\), respectively.

Let \( N_1 \) and \( N_2 \) be the two points to be connected with a cable. Let \( e_1, e_2 \) denote the slope and the intercept of the straight line \( N_1N_2 \). The coordinates of \( N_1 \) and \( N_2 \) are \((x_{N_1}, y_{N_1})\) and \((x_{N_2}, y_{N_2})\), where \( x_{N_1} < x_{N_2}, y_{N_1} = e_1x_{N_1} + e_2 \) and \( y_{N_2} = e_1x_{N_2} + e_2 \). The distance between \( N_1 \) and \( N_2 \) is

\[
\hat{d} = |x_{N_2} - x_{N_1}| \sqrt{1 + e_1^2}.
\]

A cable is defined as a continuous curve that connects \( N_1 \) and \( N_2 \) in \( \mathbb{R}^2 \). Let \( \mathcal{L} \) be the set of points on the cable. Let \( \hat{L} \) be the length of the cable and \( C \) be the cost of the cable. The cable cost is assumed to be a linear function of its length. Thus, \( C = c\hat{L} \), where \( c \) is the cost per unit length.

![Fig. 1. The locations of the two points and the EFL.](image)

We further assume that

\[
\frac{\hat{D}}{2} < e_2 < \frac{\hat{D}}{2},
\]

\[
y_{N_1}y_{N_2} < 0,
\]

\[
\hat{D} \gg \hat{d}.
\]

These assumptions are made to ensure that the two points of the cable are located on each side of the EFL, and the length of the EFL is sufficiently longer than the distance between the two points so that laying a cable between the two points without crossing the fault line is too costly to be considered as the optimal solution. This is a reasonable assumption because the length of some earthquake faults is many hundreds of kilometers and crossing them is unavoidable. An example of the EFL and the two points is illustrated in Fig. 1.

In this paper, we assume that the distribution of an earthquake effect is described by a series of concentric elliptical isoseismals, whose major axes are on the EFL. They have a common eccentricity denoted by \( e \), which is assumed to be fixed for every possible epicenter on the EFL. Therefore, throughout the paper, all the ellipses we analyze and discuss
share a common eccentricity $e$. The equation of an ellipse given that an epicenter occurs at $S$ is

$$\frac{(x-x_S)^2}{a^2} + \frac{y^2}{b^2} = 1,$$

where $a$ and $b$ are the semi-major axis and semi-minor axis, and $b^2 = (1-e^2)a^2 = e^2a^2$. For historical earthquake events, the author of [51] proposed a method of drawing an individual elliptical isoseismal to enclose all observations of a particular intensity or ground motion measure through regression analysis. Then the semi-axes $a, b$ and the eccentricity $e$ can be obtained by the derived elliptical isoseismal. Fig. 2 illustrates an earthquake with elliptical isoseismals. For every point on an isolated isoseismal, the earthquake effect is same.

![Fig. 2. A visual description of an elliptical earthquake.](image)

For any point $S \in D$ and a point $M \in \mathbb{R}^2$, we define the equivalent distance from $S$ to $M$ as follows.

$$d(M, S) = \sqrt{(M - S)^T Q(M - S)},$$

where $Q = \begin{bmatrix} 1 & 0 \\ 0 & 1/e^2 \end{bmatrix}$. In fact, the equivalent distance $d(M, S)$ is the length of the semi-major axis of the elliptical isoseismal (whose epicenter is $S$ and eccentricity is $e$) on which the point $M$ is located. In Fig. 2, $P_1$ and $P_2$ share the same equivalent distance to $S$ which is $|SP_1|$. Furthermore, the equivalent distance from $S$ to the cable is defined as

$$d(S) = \min_{M \in \mathcal{L}} d(M, S),$$

i.e., $d(S)$ is equal to the length of the shortest semi-major axis of the ellipses (whose epicenter is $S$ and eccentricity is $e$) that intersect with the cable. Because strong motions and effects of an earthquake tend to attenuate as the distance increases between the earthquake epicenter and the locations considered, we assume that all points on a given ellipse have the same earthquake effect on the cable, and that a point closer to the epicenter in terms of the equivalent distance than another implies stronger motions and higher cable break probability. Therefore, the shorter the equivalent distance is from the epicenter to the cable, the more likely the cable is to break in a case of an earthquake. Accordingly, we assume that the conditional cable break (CCB) probability function, given that an earthquake epicenter occurs at $S$ and the eccentricity of the elliptical isoseismals is $e$, is a monotonically decreasing derivable function of $d(S)$, and denoted by,

$$P(\text{break} \mid S) = q(d(S)).$$

Throughout the paper, we use $q(\cdot)$ as a generalized CCB function to derive convenient expressions obtained by probabilistic analysis. Let $p(S)$ be the continuous probability density function of an earthquake epicenter occurring at $S$ on $D$. The CBB is therefore given by

$$P(\text{break}) = \int_{S \in D} P(\text{break} \mid S)p(S)dS. \quad (4)$$

**Remark.** In cases where we know the location of the fault line but we have very limited information on the geology and the history of earthquakes and ground motion, then it is appropriate to assume an earthquake epicenter is uniformly distributed on the EFL [52], [53], i.e., $p(S)$ is a uniform distribution probability density function. However, in cases where extensive information is available, e.g. by USGS in California [54], more accurate prediction on the location of the epicenter may be made using other distributions. The results in the paper are not limited to a specific probability density function of epicenter’s location. In addition, any non-negative decreasing derivable function can be used as the CCB function $q(\cdot)$, although we will adopt exponential decay function as an example in the numerical illustration later.

### III. Description and Intuitive Justification of the Cable Shape Alternatives

The straight line is the least cost option of laying the cable. The insight that will be gained from Section IV points out to the benefit of positioning the cable vertically to the EFL which intuitively reduces the length of the cable in the “danger zone” that may be affected by an earthquake. The term “danger zone” is loosely defined to mean the area close to the EFL where the probability of cable break is significant. Note that the earthquake epicenter can occur anywhere on the EFL and cable breaks occur according to the model we have described. If the CCB function is characterized by, for example, an exponential distribution, and the probability density function of the earthquake epicenter is characterized by uniform distribution, then it is difficult to assess where exactly the probability of cable break is “significant” and where it is not. Despite the fact that the “danger zone” definition is somewhat loose, it is appropriate for intuitive explanations.

![Fig. 3. A visual description of TCS.](image)

Generally speaking, and ignoring third dimensional considerations such as the earth curvature or the topography, laying a cable in a straight line or in a sequence of a few straight line segments is a common industry practice [55]. Observing the map in [56], for example, it is clear that the shape of
cables, even if they are of length of thousands of kilometers, is a sequence of few straight lines. When there is no risk or third dimensional considerations, laying a cable in a straight line is cheaper and simpler. Accordingly, in our case, no more than three straight line segments are considered between $N_1$ and $N_2$: two lines that one of their endpoints is either $N_1$ and $N_2$ and a third line that connects these two lines and crosses the EFL through the “danger zone”. We call such a cable shape that comprises three segments: Three Connected Segments (TCS). Based on the insight that will be gained from Section IV, having the middle segment vertical to the EFL will give the lowest risk in terms of CBP and this is indeed the option used in practice by cable surveyors [55]. On the other hand, a straight line between $N_1$ and $N_2$ will give the lowest cost. Accordingly, the TCS provides a wide range of practical cable shape alternatives for which tradeoffs between cost and risk can be assessed.

In Section V, we derive expressions for the CBP and cost for a TCS cable shape as a function of the length and the slope of the middle cable segment. From these expressions we can derive the Pareto front for the two objectives by varying the length and the slope of the middle segment, after we establish the monotonicity of the CBP and cable cost function.

Further significant CBP reduction, for an additional cable length cost, can be achieved by Hook with Right Angle (HRA), where we position a segment of cable parallel to the EFL and use two vertical segments to connect it to the two endpoints. Both the benefit in CBP reduction and the cable cost depends on the distance between the parallel segment and the EFL. Therefore, we vary in Section VI this distance and study its effect on cost and risk. HRA is especially applicable to the cases where the endpoints are located far from each other but their distance to the EFL is not too far. Consider the limit where their distance to the EFL is fixed, but the distance between the endpoints is arbitrarily large. Under such a limit, it is clear that the main part of the cable under HRA is the segment that is parallel to the EFL. HRA keeps the cable parallel to the EFL for a sufficiently long distance. In the limit, the cost of HRA approaches the cost of the straight line but if the distance of the parallel segment from the EFL is far enough (this has negligible cost implication in this case), HRA can achieve very low risk of cable break. Generally, HRA has lower risks than TCS which is demonstrated numerically in Section VI. The two sets (TCS and HRA) of cable alternatives cover a wide range of CBP variations from the least cost straight line case to the low risk case of HRA.

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In this section we only consider cases where the two points are connected by a single cable in a straight line. First we derive the CBP of such a cable, and then, we consider various angles between the cable and the EFL to study their effect on CBP.

A. Line segment between $N_1$ and $N_2$

Without loss of generality, we assume that the slope of $N_1 N_2$ is greater than 0. We further assume, without loss of generality, that $N_1$ is closer to the EFL than $N_2$. Thus, we have $e_1 > 0$ and $|y_{N_1}| < |y_{N_2}|$.

The straight line segment between $N_1$ and $N_2$ has important applications throughout the paper. The assumptions we have made above, that are without loss of generality, are also relevant to the various analyzes in Sections V and VI.

We restrict the cable function to be $y = e_1 x + e_2$, where $x_{N_1} \leq x \leq x_{N_2}$, as shown in Fig. 5. Therefore, for any point $S = (x, 0)$ on the EFL, we have

$$d(S) = \min_{M \in L} d(M, S)$$

$$= \min_{M \in L} \sqrt{(x_M - x)^2 + (e_1 x_M + e_2)^2 / e^2}. \quad (5)$$

Hereinafter, we will use interchangeably the notation $d(S)$ and $d(x)$ for the equivalent distance from $S$ to the cable. Let the derivative of $d(M, S)$ (of the variable $x_M$) be 0, then we obtain

$$x_* = x_M + e_1 (e_1 x_M + e_2) / e^2 = x_M + e_1 y_M / e^2.$$

Since we have $x_{N_1} \leq x_M \leq x_{N_2}$, let $x_{S_1} = x_{N_1} + e_1 y_{N_1} / e^2$ and $x_{S_2} = x_{N_2} + e_1 y_{N_2} / e^2$. Here $S_1, S_2$ are the epicenters of the two ellipses (whose eccentricity is $e$) such that $N_1 N_2$ is the tangent at $N_1$ and $N_2$ to the two ellipses respectively, as shown in Fig. 5. According to value of $x$, there are three candidates of $d(S)$, which are

$$d_1(x) = \sqrt{(x_{N_1} - x)^2 + y_{N_1}^2 / e^2}, \quad x \leq x_{S_1},$$

$$d_2(x) = |e_1 x + e_2| / \sqrt{e_1^2 + e^2}, \quad x_{S_1} < x < x_{S_2}, \quad (6)$$

$$d_3(x) = \sqrt{(x_{N_2} - x)^2 + y_{N_2}^2 / e^2}, \quad x \geq x_{S_2}.$$

The three candidates for $d(S)$ have the following geometric meanings:

![Fig. 5. A visual description of the problem with straight line shape cable](image)
1) \( d_1(x) \): the equivalent distance from \( S \) to \( N_1 \); 
2) \( d_2(x) \): the equivalent distance from \( S \) to \( N_1N_2 \); 
3) \( d_3(x) \): the equivalent distance from \( S \) to \( N_2 \).

Then, the CBP (with the assumptions \( x_A < x_{S_1} \) and \( x_B > x_{S_2} \)) is obtained by

\[
P(\text{break}) = \frac{\int_{-\hat{D}/2}^{x_{N_2} + e_1 y_{N_1}/\varepsilon^2} q(\sqrt{(x_{N_1} - x)^2 + y_{N_1}^2/\varepsilon^2}) p(x) dx}{\int_{-\hat{D}/2}^{x_{N_2} + e_1 y_{N_1}/\varepsilon^2} q(x) dx} + \frac{\int_{x_{N_1} + e_2 y_{N_1}/\varepsilon^2}^{x_{N_2} + e_1 y_{N_1}/\varepsilon^2} q(\sqrt{(x_{N_1} - x)^2 + y_{N_1}^2/\varepsilon^2}) p(x) dx}{\int_{x_{N_1} + e_2 y_{N_1}/\varepsilon^2}^{x_{N_2} + e_1 y_{N_1}/\varepsilon^2} q(x) dx} + \frac{\int_{x_{N_2} + e_1 y_{N_1}/\varepsilon^2}^{\hat{D}/2} q(\sqrt{(x_{N_2} - x)^2 + y_{N_2}^2/\varepsilon^2}) p(x) dx}{\int_{x_{N_2} + e_1 y_{N_1}/\varepsilon^2}^{\hat{D}/2} q(x) dx}.
\]

To illustrate the CBP in a numerical example, an explicit CCB function is required. In seismology, exponential distributions are commonly used for seismic hazard assessments [57]–[59]. Hence, in our numerical examples, for simplicity of exposition and without loss of generality, we consider a negative exponential function as the CCB function and a uniformly distributed probability density function for epicenters to obtain numerical results, which are

\[
P(\text{break} | S) = q(d(S)) = e^{-\lambda d(S)}, \quad p(S) = \frac{1}{\hat{D}} \quad (7)
\]

where \( \lambda \) is a given parameter representing the exponential decay rate of the conditional cable break probability. The parameter \( \lambda \) may be estimated from repair records of buried utilities during past earthquakes. For example, the 1994 Northridge earthquake in Los Angeles caused substantial damage to the water distribution system in Los Angeles. The water trunk line (nominal pipe diameter \( \geq 600\text{mm} \)) was damaged at 74 locations (e.g., [9]), and water distribution pipelines (nominal pipe diameter \( < 600\text{mm} \)) were repaired at 1013 locations (e.g., [60]). Using these repair records, Jeon and O’Rourke [60] have developed correlations between pipeline repair rate (repairs/km) and peak ground velocity (PGV) for pipelines with different types of material. Understanding of the relationship between PGV and repair rate can lead to approximations of the (decreasing) CCB function and in the particular case of the exponential function to the estimation of \( \lambda \) in the area under consideration. For the example based on the parameters described in Table I, the CBP is 0.3379 and the cable cost is 492,7038.

**TABLE I**

<table>
<thead>
<tr>
<th>((x_{N_1}, y_{N_1}))</th>
<th>((x_{N_2}, y_{N_2}))</th>
<th>(\varepsilon)</th>
<th>(\lambda)</th>
<th>(e_1)</th>
<th>(e_2)</th>
<th>(\hat{D})</th>
<th>(\varepsilon)</th>
</tr>
</thead>
<tbody>
<tr>
<td>((-216, -27))</td>
<td>((270, 54))</td>
<td>(\frac{\hat{D}}{2})</td>
<td>(\frac{1}{3})</td>
<td>(\frac{1}{6})</td>
<td>(9)</td>
<td>(1080)</td>
<td>(1)</td>
</tr>
</tbody>
</table>

**B. Rotated straight line cable around \( N_1 \)**

Let the angle between the EFL and the straight line cable be \( \alpha \), see Fig. 6. Here we rotate the straight line cable around \( N_1 \). to observe how \( \alpha \) affects CBP. We fix \( N_1 \) and keep the cable length unchanged. We rotate the cable around \( N_1 \) and increase the angle \( \alpha \) until \( \alpha = 90^\circ \). The new endpoint of the cable is denoted by \( N_2 \). We redefine \( d_2(S) \) as the equivalent distance from \( S \) to \( N_1N_2 \) and redefine \( d_3(S) \) as the equivalent distance from \( S \) to \( N_1N_2 \). Let \( \alpha \) be the slope of \( N_1N_2 \). Then, the linear function that describes the cable \( N_1N_2 \) is given by

\[
y = a(x - x_{N_1}) + y_{N_1},
\]

where \( a \geq e_1 \). Then the coordinates of \( N_\alpha \) are

\[
\begin{align*}
x_{N_\alpha} &= x_{N_1} + \frac{\hat{D}}{\sqrt{1 + a^2}}, \\
y_{N_\alpha} &= y_{N_1} + \frac{a \hat{D}}{\sqrt{1 + a^2}}.
\end{align*}
\]

For any point \( S = (x, 0) \) on the EFL, the three possible candidates for \( d(S) \) are

\[
d_1(x) = \sqrt{(x - x_{N_1})^2 + y_{N_1}^2/\varepsilon^2}, \quad d_2(x) = |a(x - x_{N_1}) + y_{N_1}|/\sqrt{a^2 + \varepsilon^2}, \quad d_3(x) = \sqrt{(x - x_{N_1})^2 + y_{N_2}^2/\varepsilon^2}.
\]

Let \( S_1, S_2 \in \mathcal{D} \) be the epicenters of the two ellipses (whose eccentricity is \( e \)) such that \( N_1N_2 \) is the tangent at \( N_1 \) and \( N_\alpha \) to the two ellipses, respectively (see Fig. 6). Note that such two ellipses always exist under the assumptions made in (1) that the distance between \( N_1 \) and \( N_2 \) is significantly smaller than the length of the EFL). We obtain that

\[
x_{S_1} = x_{N_1} + \frac{a}{\varepsilon^2} y_{N_1}, \quad x_{S_2} = x_{N_\alpha} + \frac{a}{\varepsilon^2} y_{N_\alpha}.
\]

There are four cases for the CBP:

1) **When \( x_{S_1} \geq -\hat{D}/2 \) and \( x_{S_2} \leq \hat{D}/2 \):**

\[
P(\text{break}) = \int_{-\hat{D}/2}^{x_{S_1}} q(d_1(x)) p(x) dx + \int_{x_{S_1}}^{\hat{D}/2} q(d_2(x)) p(x) dx + \int_{x_{S_2}}^{\hat{D}/2} q(d_3(x)) p(x) dx,
\]
2) When \( x_{S_1} \geq -\hat{D}/2 \) and \( x_{S_2} > \hat{D}/2 \):

\[
P(\text{break}) = \int_{-\hat{D}/2}^{x_{S_1}} q(d_1(x))p(x)dx + \int_{x_{S_1}}^{\hat{D}/2} q(d_2(x))p(x)dx,
\]

3) When \( x_{S_1} < -\hat{D}/2 \) and \( x_{S_2} \leq \hat{D}/2 \):

\[
P(\text{break}) = \int_{x_{S_2}}^{\hat{D}/2} q(d_2(x))p(x)dx + \int_{x_{S_2}}^{\hat{D}/2} q(d_3(x))p(x)dx,
\]

4) When \( x_{S_1} < -\hat{D}/2 \) and \( x_{S_2} > \hat{D}/2 \):

\[
P(\text{break}) = \int_{-\hat{D}/2}^{\hat{D}/2} q(d_2(x))p(x)dx,
\]

where \( d_1(x) - d_3(x) \) are given by (8)-(10), and \( x_{S_1}, x_{S_2} \) are given by (11), (12), respectively.

V. THREE CONNECTED SEGMENTS

The straight line cable costs the least, but its CBP is high. In this section, we consider a cable shape called Three Connected Segments (TCS), where we restrict the set of candidate cable functions to a shape described by Fig. 3. Each cable we consider consists of three segments: \( \overrightarrow{N_1D}, \overrightarrow{DE} \) and \( \overrightarrow{EN_2} \). The middle segment \( \overrightarrow{DE} \) always passes through the point \( C \), where \( C \) is the crossing point of \( N_1N_2 \) and the EFL. The coordinates of \( C \) are \((-e_2/e_1, 0)\). Let the angle between the middle segment \( \overrightarrow{DE} \) and the \( y \)-axis be \( \theta \). Let \( m \) be the slope of \( \overrightarrow{DE} \). Then the linear function that describes \( \overrightarrow{DE} \) is given by

\[
y = m(x + \frac{e_2}{e_1}),
\]

where \( m = \tan(90^\circ - \theta) \) and \( m \geq e_1 \). We define the parameter \( \Delta_T \) to be equal to length of the middle segment \( \overrightarrow{DE} \). There is a one-to-one correspondence between the value of \( \Delta_T \) and the set of cable functions under TCS. For each value of \( \Delta_T \), we ensure that the triangles \( N_1CD \) and \( N_2CE \) are two similar triangles. Thus, we have

\[
\frac{|N_1C|}{|CN_2|} = \frac{|DC|}{|CE|} = \frac{|N_1D|}{|EN_2|}.
\]

Denote the slope of \( \overrightarrow{N_1D} \) and \( \overrightarrow{N_2E} \) by \( k \). We obtain the expression of \( k \), which is

\[
k(\Delta_T, m) = e_1 + \frac{(m - e_1)\Delta_T}{\Delta_T - (x_{N_2} - x_{N_1})\sqrt{1 + m^2}}.
\]

For any point \( S = (x, 0) \) on the EFL, we have

\[
d(S) = \min_{M \in \mathcal{L}} d(M, S).
\]

We take the derivative of \( d(M, S) \) (of the variable \( x_M \)), and let it be equal to 0. Then we obtain that there are five possible candidates for \( d(S) \), which are

\[
d_{T_1}(S) = d_{T_1}(x) = \sqrt{(x - x_{N_1})^2 + y^2_{N_1}/e_1^2}, \quad (13)
\]

\[
d_{T_2}(S) = d_{T_2}(x, \Delta_T, m) = \frac{|kx - kx_{N_1} + y_{N_1}|}{\sqrt{k^2 + e_2^2}}, \quad (14)
\]

\[
d_{T_3}(S) = d_{T_3}(x, m) = \frac{|mx + me_2/e_1|}{\sqrt{m^2 + e_2^2}}, \quad (15)
\]

\[
d_{T_4}(S) = d_{T_4}(x, \Delta_T, m) = \frac{|kx - kx_{N_2} + y_{N_2}|}{\sqrt{k^2 + e_2^2}}, \quad (16)
\]

and

\[
d_{T_5}(S) = d_{T_5}(x) = \sqrt{(x - x_{N_2})^2 + y^2_{N_2}/e_2^2}. \quad (17)
\]

The five candidates for \( d(S) \) have the following geometric meanings:

1) \( d_{T_1}(S) \): the equivalent distance from \( S \) to \( N_1 \);
2) \( d_{T_2}(S) \): the equivalent distance from \( S \) to \( \overrightarrow{N_1D} \);
3) \( d_{T_3}(S) \): the equivalent distance from \( S \) to \( \overrightarrow{EN_2} \);
4) \( d_{T_4}(S) \): the equivalent distance from \( S \) to \( \overrightarrow{N_2E} \);
5) \( d_{T_5}(S) \): the equivalent distance from \( S \) to \( N_2 \).

We see that only the functions \( d_{T_2}(x, \Delta_T, m) \) and \( d_{T_4}(x, \Delta_T, m) \) depend on \( \Delta_T \). This implies that, for a given value of \( m \), the value of \( \Delta_T \) can dictate which is the smallest distance among the five candidates for \( d(S) \). Therefore, according to the value of \( \Delta_T \), the problem is further divided into two scenarios as shown in Fig. 8. Next, we find the interval of \( \Delta_T \) for each of the two scenarios.

Let \( S_{T_5} \in AC \) be the epicenter of the ellipse (whose eccentricity is \( e \)) such that \( \overrightarrow{DE} \) is the tangent to the ellipse and \( N_1 \) is on the ellipse. Similarly, let \( S_{T_6} \in CB \) be the epicenter of another ellipse (whose eccentricity is \( e \)) such that \( \overrightarrow{DE} \) is the tangent to the ellipse and \( N_2 \) is on the ellipse. Then, we have \( d(S_{T_5}) = d_{T_1}(S_{T_5}) = d_{T_3}(S_{T_5}) \) and
\[ d(S_{T6}) = d_{T3}(S_{T6}) = d_{T5}(S_{T6}). \] The x-coordinates of \( S_{T5} \) and \( S_{T6} \) are
\[ x_{S_{T5}}(m) = x_{N_1} + \left( m^2 - \sqrt{(m^2 + \varepsilon^2)(m^2 - e_1^2)} \right) \frac{y_{N_1}}{e_1\varepsilon^2}, \tag{18} \]
and
\[ x_{S_{T6}}(m) = x_{N_2} + \left( m^2 - \sqrt{(m^2 + \varepsilon^2)(m^2 - e_1^2)} \right) \frac{y_{N_2}}{e_1\varepsilon^2}, \tag{19} \]
respectively. There is a value of \( \Delta_T \) such that \( N_1D \) is the tangent at \( N_1 \) to the ellipse with the epicenter \( S_{T5} \), and \( N_2E \) is the tangent at \( N_2 \) to the ellipse with epicenter \( S_{T6} \), respectively. We obtain this boundary value of \( \Delta_T \), which is
\[ \Delta_{T\text{-bound}}(m) = \frac{\sqrt{1 + m^2(x_{N_2} - x_{N_1})(m^2 - e_1^2) - \sqrt{(m^2 + \varepsilon^2)(m^2 - e_1^2)}}}{m^2 - e_1m - \sqrt{(m^2 + \varepsilon^2)(m^2 - e_1^2)}}. \tag{20} \]

The scenario described in Fig. 8(a) corresponds to the case \( 0 \leq \Delta_T \leq \Delta_{T\text{-bound}}(m) \). For this scenario, let \( S_{T1} \in \overline{AC} \) be the epicenter of the ellipse (whose eccentricity is \( e \)) such that \( N_1D \) is the tangent at \( N_1 \) to the ellipse. Similarly, let \( S_{T4} \in \overline{CB} \) be the epicenter of another ellipse (whose eccentricity is \( e \)) such that \( N_2E \) is the tangent at \( N_2 \) to the ellipse. Then, we have
\[ d(S_{T1}) = d_{T1}(S_{T1}) = d_{T2}(S_{T1}) \] and \( d(S_{T4}) = d_{T3}(S_{T4}) = d_{T5}(S_{T4}) \). Let \( S_{T2} \in \overline{ST_1} \) be the epicenter of the ellipse (whose eccentricity is \( e \)) such that \( DE \) is the tangent to the ellipse and \( N_1D \) is the tangent at \( E \) to the ellipse, where \( F \in \overline{N_1D} \). Similarly, let \( S_{T3} \in \overline{ST_4} \) be the epicenter of another ellipse (whose eccentricity is \( e \)) such that \( DE \) is the tangent to the ellipse and \( N_2E \) is the tangent at \( G \) to the ellipse, where \( G \in \overline{N_2E} \). Then, we have
\[ d(S_{T2}) = d_{T2}(S_{T2}) = d_{T3}(S_{T2}) \] and \( d(S_{T3}) = d_{T3}(S_{T3}) = d_{T4}(S_{T3}) \). The x-coordinates of these four points are
\[ x_{S_{T1}}(\Delta_T, m) = x_{N_1} + ky_{N_1}/\varepsilon^2, \tag{21} \]
\[ x_{S_{T2}}(\Delta_T, m) = \frac{(kx_{N_1} - y_{N_1})\sqrt{m^2 + \varepsilon^2 + m \varepsilon^2 \sqrt{m^2 + \varepsilon^2}}}{m \sqrt{m^2 + \varepsilon^2 - m \sqrt{m^2 + \varepsilon^2}}}, \tag{22} \]
\[ x_{S_{T3}}(\Delta_T, m) = \frac{(kx_{N_2} - y_{N_2})\sqrt{m^2 + \varepsilon^2 + m \varepsilon^2 \sqrt{m^2 + \varepsilon^2}}}{m \sqrt{m^2 + \varepsilon^2 - m \sqrt{m^2 + \varepsilon^2}}}, \tag{23} \]
and
\[ x_{S_{T4}}(\Delta_T, m) = x_{N_2} + ky_{N_2}/\varepsilon^2, \tag{24} \]
respectively. In this case, we have
\[ d(S) = \begin{cases} d_{T1}(S), & S \in \overline{AS_{T1}}, \\ d_{T2}(S), & S \in \overline{ST_1ST_2}, \\ d_{T3}(S), & S \in \overline{ST_2ST_3}, \\ d_{T4}(S), & S \in \overline{ST_3ST_4}, \\ d_{T5}(S), & S \in \overline{ST_4B}. \end{cases} \]

The CBP is obtained by
\[ P(\text{break}) = \int_{-\Delta/2}^{\Delta/2} q(d_{T1}(x))p(x)dx + \int_{x_{S_{T1}}}^{x_{S_{T2}}} q(d_{T2}(x, \Delta_T, m))p(x)dx + \int_{x_{S_{T1}}}^{x_{S_{T2}}} q(d_{T3}(x, m))p(x)dx + \int_{x_{S_{T3}}}^{x_{S_{T4}}} q(d_{T4}(x, \Delta_T, m))p(x)dx + \int_{x_{S_{T4}}}^{\Delta/2} q(d_{T5}(x))p(x)dx, \tag{25} \]
where \( d_{T1}(x) - d_{T4}(x) \) are given by (13) - (17), and \( x_{S_{T1}}, x_{S_{T4}} \) are given by (21) - (24).

**Proposition 1.** The function \( P(\text{break}) \) given by (25) is monotonically non-increasing in \( \Delta_T \) for \( \Delta_T \in (0, \Delta_{T\text{-bound}}(m)) \).

**Proof:** See the Appendix. \[ \square \]

The scenario described in Fig. 8(b) corresponds to the case \( \Delta_T \geq \Delta_{T\text{-bound}}(m) \). For this scenario, we have
\[ d(S) = \begin{cases} d_{T1}(S), & S \in \overline{AS_{T5}}, \\ d_{T3}(S), & S \in \overline{ST_5ST_6}, \\ d_{T5}(S), & S \in \overline{ST_6B}. \end{cases} \]

The CBP is obtained by
\[ P(\text{break}) = \int_{-\Delta/2}^{\Delta/2} q(d_{T1}(x))p(x)dx + \int_{x_{S_{T5}}}^{x_{S_{T6}}} q(d_{T3}(x, m))p(x)dx + \int_{x_{S_{T6}}}^{\Delta/2} q(d_{T5}(x))p(x)dx, \tag{26} \]
where \( d_{T1}(x), d_{T3}(x, m), d_{T5}(x) \) are given by (13), (15), (17), and \( x_{S_{T5}}, x_{S_{T6}} \) are given by (18), (19), respectively. The CBP given by (26) is independent of \( \Delta_T \).
In the first scenario, the CBP given by (25) does not increase as $\Delta_T$ increases. In the second scenario, the CBP given by (26) is independent of $\Delta_T$. Hence, for a given value of $m$, the CBP is monotonically non-increasing in $\Delta_T$ reaching its minimum for $\Delta_T \geq \Delta_{T, \text{bound}}(m)$.

Intuitively, when increasing the length of the middle segment, the length of the entire cable increases. The cost of the cable is obtained by

$$C(\Delta_T, m) = c \left( x_{N_2} - x_{N_1} - \frac{\Delta_T}{\sqrt{1 + m^2}} \right) \sqrt{1 + k^2} + c\Delta_T,$$

which is monotonically increasing in $\Delta_T$.

Fig. 9. A special case of TCS ($\theta = 0^\circ$).

Fig. 9 illustrates a special case of TCS, where the middle cable segment is positioned vertically to the EFL. In this case, $\theta = 0^\circ$ and $m = \infty$.

In summary, we list the procedure of our method to generate the Pareto front under TCS shape as follows.

- **Input**: Two endpoints of the EFL, $A$ and $B$; two points to be connected with a cable, $N_1$, $N_2$; cost of the cable per unit length, $c$; eccentricity $e$ of the elliptical isoseisms; probability density function of location of epicenters; CCB function;
- For each $\theta \in [0, 90^\circ - \arctan e]$,
  - For each $\Delta_T > 0$,
    * Calculate CBP by (25) or (26);
    * Calculate the cable cost by (27);
- Generate Pareto front curves by each pair of the CBP and the cable cost corresponding to the same $\Delta_T$.

To illustrate the monotonicity of the CBP, we use (7) as the explicit CCB function and the spatial distribution of epicenters. We again use the parameters of Table I, and verify the expressions (25) and (26) by simulations which are shown in Fig. 10. For each of the curves, the CBP decreases and reaches a plateau, as $\Delta_T$ increases. The analytical results are always within the confidence intervals of the simulation results.

Having established the monotonicity of both the CBP and the cable cost functions, we generate the Pareto front (see Fig. 11) for these two objectives of the cables under TCS using the parameters of Table I. In each of the curves the angle is fixed as indicated. Then the curves in the figure are obtained by varying the length of the middle segment for each curve. The vertical dashed lines correspond to a range of costs for which the CBP does not change, and therefore they do not belong to the Pareto front curves. With the Pareto front, we also have solutions for the optimization problems of minimizing cost subject to a constraint on the CBP and of minimizing CBP subject to a constraint on the cable cost. In Fig. 11, all the curves meet at the point where the middle segment is of length equal to zero. Then they all represent a straight line between $N_1$ and $N_2$. At that point we obtain the minimal cost to be equal to 492.7038 and the CBP to be equal to 0.3379. This is the case of the straight line cable, which we studied in Section IV-A.

We observe that for every $\theta$ value there is a minimum CBP that can be reached by extending the length of the middle cable. Then further reductions in CBP are not achievable because the further extension is already out of the “danger zone”. We also observe that lower $\theta$ values give lower minimum CBP values, where the case $\theta = 0^\circ$ gives the lowest minimum CBP value. This is consistent with the examples presented in Section IV, and the intuitive explanation that the vertical line has minimal exposure to the “danger zone”. It is also consistent with practice used by cable surveyors [55].

Fig. 10. CBP versus $\Delta_T$ for the cables under TCS.

Fig. 11. Pareto Front for CBP and Cost of the cables under TCS.
VI. HOOK WITH RIGHT ANGLES

Based on the insight gained from the cable rotation exercise in Section IV, to reduce the CBP, another option is to position the cable perpendicular to the EFL to minimize the number of points on the cable that may be affected by an earthquake, and continue in a straight line to move away from the fault as fast as possible until reaching a safe ground on the other side. Then, lay the cable in parallel to the EFL for a sufficiently long distance and finally return to the destination. The effectiveness of this approach depends on the length it is laid on safe ground in parallel to the EFL. Recall that we assume without loss of generality that \( N_1 \) is closer to the EFL than \( N_2 \). Otherwise, this HRA cables shape will start at \( N_2 \) and end at \( N_1 \).

This gives rise to another three-segment cable shape called Hook with Right Angles (HRA), where we restrict the set of candidate cable functions to a shape described by Fig. 4. Each cable we consider consists of three segments: \( N_1D, DE \) and \( EN_2 \). We position the first cable segment \( N_1D \) perpendicular to the EFL. The crossing point of \( N_1D \) and the EFL is denoted by \( C \). We position the middle segment \( DE \) parallel to the EFL and position the last cable segment \( EN_2 \) perpendicular to \( DE \). The shape of the entire cable resembles a hook with three segments and two right angles. We define the parameter \( \Delta_H \) to be equal to the length of the cable segment \( EN_2 \). There is a one-to-one correspondence between the value of \( \Delta_H \) and the set of cable functions under HRA.

For any point \( S = (x, 0) \) on the EFL, we have

\[
d(S) = \min_{M \in L} d(M, S).
\]

We take the derivative of \( d(M, S) \) (of the variable \( x_M \)), and let it be equal to 0. Then we obtain that there are three possible candidates for \( d(S) \), which are

\[
d_{H1}(S) = d_{H1}(x) = |x - x_{N_1}|, \quad d_{H2}(S) = d_{H2}(\Delta_H) = (y_{N_2} + \Delta_H)/\varepsilon, \quad d_{H3}(S) = d_{H3}(x) = \sqrt{(x - x_{N_2})^2 + y_{N_2}^2}/\varepsilon. \tag{28, 29, 30}
\]

The three candidates for \( d(S) \) have the following geometric meanings:

1) \( d_{H1}(S) \): the equivalent distance from \( S \) to \( C \);
2) \( d_{H2}(S) \): the equivalent distance from \( S \) to \( DE \);
3) \( d_{H3}(S) \): the equivalent distance from \( S \) to \( N_2 \).

Let \( F \) be orthogonal projection of \( N_2 \) on the EFL. When \( |CF| \leq |FN_2|/\varepsilon \), for any \( S \in CF \), the equivalent distance \( d(S) \) is always \( d_{H1}(S) \). In this scenario, the CBP is not a function of \( \Delta_H \). Increasing \( \Delta_H \) does not help in reducing the CBP. However, in the scenario that \( |CF| > |FN_2|/\varepsilon \), increasing \( \Delta_H \) reduces the CBP. Therefore, in the following, we only consider the scenario that \( |CF| > |FN_2|/\varepsilon \), which is in particular

\[
x_{N_2} - x_{N_1} > y_{N_2}/\varepsilon. \tag{31}
\]

Equation (31) ensures the effectiveness of HRA, and as mentioned above, it is related to having sufficient length parallel to the EFL in “safe ground”.

With a fixed value of \( \Delta_H \), the function \( d_{H2}(\Delta_H) \) is constant for any \( S \in D \). And we see that only \( d_{H2}(\Delta_H) \) depends on \( \Delta_H \). This implies that the value of \( \Delta_H \) can dictate which is the smallest distance among the three candidates for \( d(S) \). Therefore, according to the value of \( \Delta_H \), the problem is further divided into two scenarios as shown in Fig. 12.

\[
\begin{align*}
\Delta_H &> \frac{\varepsilon(x_{N_2} - x_{N_1})^2 + y_{N_2}^2/\varepsilon}{2(x_{N_2} - x_{N_1}) - y_{N_2}}. \tag{32}
\end{align*}
\]

The scenario described in Fig. 12(a) corresponds to the case \( \Delta_H \geq \frac{d_{H1}(S_{H0})}{d_{H3}(S_{H0})} \). The range of \( \Delta_H \) for this scenario is

\[
\Delta_H \geq \frac{x_{S_{H0}} - x_{N_2}}{2(y_{N_2} - y_{N_1})} - y_{N_2}.
\]

In this case, we have

\[
d(S) = \begin{cases} d_{H1}(S), & S \in \overline{AS_{H0}}, \\ d_{H3}(S), & S \in \overline{S_{H0}B}. \end{cases}
\]

The CBP is obtained by

\[
P(\text{break}) = \int_{-D/2}^{x_{S_{H0}}} q(d_{H1}(x))p(x)dx + \int_{x_{S_{H0}}}^{D/2} q(d_{H3}(x))p(x)dx,
\]

where \( d_{H1}(x), d_{H3}(x) \) are given by (28), (30), and \( x_{S_{H0}} \) is given by (32). The CBP in (33) is independent of \( \Delta_H \).

The scenario described in Fig. 12(b) corresponds to the case \( d_{H2}(\Delta_H) < d_{H1}(S_{H0}) = d_{H3}(S_{H0}) \). The range of \( \Delta_H \) for
this scenario is
\[ 0 \leq \Delta_H < \frac{\varepsilon(x_{N_2} - x_{N_1})^2 + y_{N_2}^2/\varepsilon}{2(x_{N_2} - x_{N_1})} - y_{N_2}. \]

For this scenario, let \( S_{H1} \in C_{SH1} \) be the epicenter of the ellipse (whose eccentricity is \( e \)) such that the point \( C \) is on the ellipse and \( DE \) is the tangent to the ellipse. Similarly, let \( S_{H2} \in C_{SH2} \) be the epicenter of another ellipse (whose eccentricity is \( e \)) such that the point \( N_2 \) is on the ellipse and \( DE \) is the tangent to the ellipse. Thus, we have \( d_{H1}(S_{H1}) = d_{H1}(\Delta_H) = d_{H2}(S_{H2}) \). The x-coordinates of \( S_{H1} \) and \( S_{H2} \) are
\[ x_{S_{H1}}(\Delta_H) = x_{N_1} + (y_{N_2} + \Delta_H)/\varepsilon, \quad (34) \]
and
\[ x_{S_{H2}}(\Delta_H) = x_{N_2} - \frac{1}{\varepsilon} \sqrt{\Delta_H^2 + 2y_{N_2}\Delta_H}, \quad (35) \]
respectively. In this case, we have
\[ d(S) = \begin{cases} d_{H1}(S), & S \in A_{S_{H1}}, \\ d_{H2}(\Delta_H), & S \in S_{H1}S_{H2}, \\ d_{H3}(S), & S \in S_{H2}B. \end{cases} \]
The CBP is obtained by
\[ P(\text{break}) = \int_{-D/2}^{x_{S_{H1}}} q(d_{H1}(x))p(x)dx + \int_{x_{S_{H1}}}^{x_{S_{H2}}} p(x)dx + \int_{x_{S_{H2}}}^{D/2} q(d_{H3}(x))p(x)dx, \quad (36) \]
where \( d_{H1}(x) - d_{H3}(x) \) are given by (28) - (30), and \( x_{S_{H1}}, x_{S_{H2}} \) are given by (34), (35).

**Proposition 2.** The function \( P(\text{break}) \) given by (36) is monotonically non-increasing in \( \Delta_H \) for
\[ \Delta_H \in \left( 0, \frac{\varepsilon(x_{N_2} - x_{N_1})^2 + y_{N_2}^2/\varepsilon}{2(x_{N_2} - x_{N_1})} - y_{N_2} \right). \]

**Proof:** See the Appendix.

In the first scenario, the CBP given by (33) is independent of \( \Delta_H \). In the second scenario, the CBP given by (36) does not increase as \( \Delta_H \) increases. Hence, for the entire range of \( \Delta_H \), the CBP is monotonically non-increasing in \( \Delta_H \) reaching its minimum for
\[ \Delta_H \geq \frac{\varepsilon(x_{N_2} - x_{N_1})^2 + y_{N_2}^2/\varepsilon}{2(x_{N_2} - x_{N_1})} - y_{N_2}. \]

The cost of a HRA cable is obtained by
\[ C(\Delta_H) = c(y_{N_2} - y_{N_1} + x_{N_2} - x_{N_1} + 2\Delta_H), \quad (37) \]
which is monotonically increasing in \( \Delta_H \).

We also summarize the procedure of our method to generate the Pareto front under HRA shape as follows.

- **Input:** Two endpoints of the EFL, \( A \) and \( B \); two points to be connected with a cable, \( N_1, N_2 \); cost of the cable per unit length, \( c \); eccentricity \( e \) of the elliptical isoseisms;

- **Probability density function of location of epicenters:**

- **CCB function:**

- **For each \( \Delta_H > 0 \),**
  
  - Calculate CBP by (33) or (36);
  
  - Calculate the cable cost by (37);

- **Generate Pareto front curves by each pair of the CBP and the cable cost corresponding to the same \( \Delta_H \).**

Finally, we would like to discuss the special cases illustrated in Fig. 14, where the middle segment of a cable still has right angles with \( N_1D \) and \( FN_2 \), but the cable does not have a hook shape, because its middle segment is between the EFL and \( N_2 \). For consistency with the previous \( \Delta_H \) definition, we define \( \Delta_H \) as taking negative values in such cases, and in particular, we consider the \( \Delta_H \) values in the range \(-|FN_2| \leq \Delta_H < 0\).

In such cases, the cable cost does not vary with \( \Delta_H \), but the CBP increases as the middle segment approaches the EFL, and reaches its maximum value when \( \Delta_H = -|FN_2| \).
maximum value is obtained by
\[
P(\text{break}) = \int_{-\frac{D}{2}}^{x_{N1}} q(|x - x_{N1}|) p(x)dx + q(0) \int_{x_{N1}}^{x_{N2}} p(x)dx + \int_{x_{N2}}^{\frac{D}{2}} q(|x - x_{N2}|) p(x)dx.
\]

Using (7) and the parameters of Table I, we obtain the numerical value of this maximum CBP, which is 0.5495.

Having established the monotonicity of both the CBP and the cable cost functions, we generate the Pareto front for these two objectives of the cables under HRA and compare it with the Pareto front of the cables under TCS (vertical middle segment case). The two Pareto fronts are shown in Fig. 15. The vertical dashed lines correspond to a range of costs for which the CBP does not change. The horizontal dashed line corresponds to a range of CBP for which the costs does not change. Therefore, the dashed lines do not belong to the Pareto front curves, based on the Pareto front definition, in which it is impossible to make one objective better off without making the other objective worse off. The Pareto front curves enable us to solve the optimization problems of minimizing cost subject to a constraint on the CBP and of minimizing CBP subject to a constraint on the cable cost.

Although our approach relies on a 2D simplification, it provides a step-stool for the more complex 3D multi-objective optimization problem, where the geography as well as the real earthquake effects are considered, and Pareto front solutions are derived.

**Appendix**

**Proof of Proposition 1:** We have
\[
\frac{dP(\text{break})}{d\Delta T} = q(d_{T1}(x_{S_{T1}})) \frac{dx_{S_{T1}}}{d\Delta T} p(x_{S_{T1}}) + q(d_{T2}(x_{S_{T2}}, \Delta_T, m)) \frac{dx_{S_{T2}}}{d\Delta T} p(x_{S_{T2}}) + q(d_{T3}(x_{S_{T3}}, \Delta_T, m)) \frac{dx_{S_{T3}}}{d\Delta T} p(x_{S_{T3}}) + q(d_{T4}(x_{S_{T4}}, \Delta_T, m)) \frac{dx_{S_{T4}}}{d\Delta T} p(x_{S_{T4}}) + q(d_{T5}(x_{S_{T5}})) \frac{dx_{S_{T5}}}{d\Delta T} p(x_{S_{T5}}) + \int_{x_{S_{T1}}}^{x_{S_{T2}}} \frac{\partial q(d_{T2}(x, \Delta_T, m))}{\partial \Delta T} p(x)dx + \int_{x_{S_{T3}}}^{x_{S_{T4}}} \frac{\partial q(d_{T4}(x, \Delta_T, m))}{\partial \Delta T} p(x)dx.
\]

In the scenario described in Fig. 8(a), we have
\[
d_{T1}(x_{S_{T1}}) = d_{T2}(x_{S_{T1}}, \Delta_T, m),
\]
\[
d_{T2}(x_{S_{T2}}, \Delta_T, m) = d_{T3}(x_{S_{T2}}, m),
\]
\[
d_{T3}(x_{S_{T3}}, m) = d_{T4}(x_{S_{T3}}, \Delta_T, m),
\]

and
\[
d_{T4}(x_{S_{T4}}, \Delta_T, m) = d_{T5}(x_{S_{T4}}).
\]

**VII. Conclusion**

We have considered the problem of how to lay a single cable to connect two points across an EFL. We have formulated the problem as a multi-objective optimization on the plane that considers cable cost and CBP as two objectives. We have gained insight from the rotation exercise of a straight-line cable and further considered two important and meaningful sets of cable shape alternatives. For each set of alternatives, we have identified the key attributes which affect both the cable cost and the CBP. We have derived the expressions that have led to the Pareto fronts for the two objectives under TCS and HRA, which enable us to solve two single objective optimization problems of minimizing cost subject to a constraint on the CBP, and of minimizing CBP subject to a constraint on the cable cost.
Thus, we have
\[
\frac{dP_{\text{break}}}{d\Delta_T} = \int_{x_S}^{x_{T_2}} \frac{\partial q_{\text{break}}(x, \Delta_T, m)}{\partial \Delta_T} \cdot p(x) dx
\]
\[
+ \int_{x_{S_T_1}}^{x_{T_2}} \frac{\partial q_{\text{break}}(x, \Delta_T, m)}{\partial \Delta_T} dx
\]
\[
= \int_{x_{S_T_1}}^{x_{T_2}} \frac{q'(d_{T_2}(x, \Delta_T, m))}{\partial \Delta_T} \cdot p(x) dx
\]
\[
+ \int_{x_{S_T_1}}^{x_{T_2}} \frac{q'(d_{T_4}(x, \Delta_T, m))}{\partial \Delta_T} \cdot p(x) dx.
\]

Since \( q'(\cdot) \) is defined as a monotonically decreasing derivable function, we have \( q'(d_{T_2}(x, \Delta_T, m)) < 0 \) and \( q'(d_{T_4}(x, \Delta_T, m)) < 0 \). We further obtain that
\[
\frac{\partial d_{T_2}(x, \Delta_T, m)}{\partial \Delta_T} = \frac{\varepsilon^2(x - x_{S_T_1})(m - 1)(x_N - x_N_1)\sqrt{1 + m^2}}{(\Delta_T - (x_N - x_N_1)\sqrt{1 + m^2})^2} (k^2 + \varepsilon^2)^{3/2},
\]
and
\[
\frac{\partial d_{T_4}(x, \Delta_T, m)}{\partial \Delta_T} = \frac{\varepsilon^2(x_{S_T_4} - x)(m - 1)(x_N - x_N_1)\sqrt{1 + m^2}}{(\Delta_T - (x_N - x_N_1)\sqrt{1 + m^2})^2} (k^2 + \varepsilon^2)^{3/2}.
\]

Note that we defined that \( m > e_1 \). Thus, for \( x \in (x_{S_T_1}, x_{T_2}) \), we have
\[
\frac{\partial d_{T_2}(x, \Delta_T, m)}{\partial \Delta_T} > 0.
\]
Similarly, for \( x \in (x_{S_T_2}, x_{S_T_4}) \), we have
\[
\frac{\partial d_{T_4}(x, \Delta_T, m)}{\partial \Delta_T} > 0.
\]

Since \( p(x) \geq 0 \), we have
\[
\frac{dP_{\text{break}}}{d\Delta_T} \leq 0.
\]

Therefore, the function \( P_{\text{break}} \) given by (25) is monotonically non-increasing in \( \Delta_T \) for \( \Delta_T \in (0, \Delta_{T,\text{bound}}(m)) \).

---

**Proof of Proposition 2:** We have
\[
\frac{dP_{\text{break}}}{d\Delta_H} = q(d_{H_1}(x_{S_{H_1}})) \cdot \frac{dx_{S_{H_1}}}{d\Delta_H} \cdot p(x_{S_{H_1}})
\]
\[
+ q(d_{H_2}(\Delta_H)) \cdot \frac{dx_{S_{H_2}}}{d\Delta_H} \cdot p(x_{S_{H_1}}) - \frac{dx_{S_{H_1}}}{d\Delta_H} \cdot p(x_{S_{H_1}})
\]
\[
+ \frac{dx_{S_{H_2}}}{d\Delta_H} \int_{x_{S_{H_1}}}^{x_{S_{H_2}}} p(x) dx.
\]
\[
- q(d_{H_3}(x_{S_{H_2}})) \cdot \frac{dx_{S_{H_2}}}{d\Delta_H} \cdot p(x_{S_{H_2}}).
\]

In the scenario described in Fig. 12(b), we have \( d_{H_1}(x_{S_{H_1}}) = d_{H_2}(\Delta_H) = d_{H_3}(x_{S_{H_2}}) \). Thus, we have
\[
\frac{dP_{\text{break}}}{d\Delta_H} = q(d_{H_2}(\Delta_H)) \int_{x_{S_{H_1}}}^{x_{S_{H_2}}} p(x) dx.
\]
\[
= \frac{1}{\varepsilon} q'(d_{H_2}(\Delta_H)) \int_{x_{S_{H_1}}}^{x_{S_{H_2}}} p(x) dx.
\]

Since \( q'(\cdot) \) is defined as a monotonically decreasing derivable function, we have \( q'(d_{H_2}(\Delta_H)) < 0 \). Furthermore, \( \int_{x_{S_{H_1}}}^{x_{S_{H_2}}} p(x) dx \geq 0 \). Thus we have
\[
\frac{dP_{\text{break}}}{d\Delta_H} \leq 0.
\]

Therefore, the function \( P_{\text{break}} \) given by (36) is monotonically non-increasing in \( \Delta_H \) for
\[
\Delta_H \in \left(0, \frac{\varepsilon(x_N - x_{N_1})^2 + y_{N_2}^2/\varepsilon}{2(x_N - x_{N_1})} - y_{N_2}\right).
\]

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